

Unit II – Divide and Conquer

- Merge sort
- Quick sort
- Binary search
- **Multiplication of large Integers**
- **Strassen's Matrix Multiplication**
 - Decrease the no. of multiplications at the expense of a slight increase in no. of additions

Multiplication of large Integers

DATE

Example

23

14

$a_0 a_1$

$b_0 b_1$

$$n = 2$$

$$n^2 = 4 \text{ Multi}$$

$$23 = (2 \times 10^1) + (3 \times 10^0)$$

$$14 = (1 \times 10^1) + (4 \times 10^0)$$

$$23 \times 14$$

$$[(2 \times 10^1) + (3 \times 10^0)] \times [(1 \times 10^1) + (4 \times 10^0)]$$

$$(2 \times 1) 10^2 + \underbrace{[(2 \times 4) + (3 \times 1)]}_{8+3=11} 10^1 + (3 \times 4) 10^0$$

$$\begin{array}{ccccccc} (2 \times 4) & + & (3 \times 1) & = & (2+3) \times (4+1) & - & 2 \times 1 - \\ a_0 & b_1 & a_1 & b_0 & & & 3 \times 4 \end{array}$$

$$= 5 \times 5 - (2 + 12)$$

$$= 25 - 14 = 11$$

Multiplication of large Integers

$$c = a * b = \underline{\underline{c_2}} 10^2 + \underline{\underline{c_1}} 10^1 + \underline{\underline{c_0}}$$

$$c_2 = 2 * 1 \Rightarrow a_0 * b_0$$

$$c_0 = 3 * 4 \Rightarrow a_1 * b_1$$

$$c_1 = (a_0 + a_1) * (b_0 + b_1) - (c_2 + c_0)$$

$$c = a * b = c_2 10^{n/2} + c_1 10^{n/2} + c_0$$

$$M(n) = 3 M(n/2) \quad \text{for } n > 1$$

$$M(1) = 1$$

$$\boxed{n = 2^k} \Rightarrow k = \log_2 n$$

$$M(2^k) = 3 M(2^k/2^1)$$

$$= 3 M(2^{k-1})$$

$$= 3 [3 M(2^{k-2})]$$

$$= 3^2 M(2^{k-2})$$

$$= 3^3 M(2^{k-3})$$

$$= 3^4 M(2^{k-4})$$

$$= 3^5 M(2^{k-5}) \Rightarrow 3^i M(2^{k-i})$$

$$\Rightarrow 3^k M(2^{k-k})$$

$$\Rightarrow 3^k M(2^0)$$

$$\Rightarrow 3^k M(1)$$

$$\Rightarrow 3^k //$$

$$\Rightarrow 3^{\log_2 n}$$

$$\Rightarrow n^{\log_2 3}$$

$$\Rightarrow \underline{\underline{n^{1.585}}}$$

$$\boxed{\frac{\log_b c}{a} = \frac{\log_b a}{c}}$$

performed well on ^{integers} 600 digit long

Multiplication of large Integers

- Example;
- $1234 * 1234$
- $123456 * 654321$

Strassen's Matrix Multiplication

- Strassen - 1969

$$\begin{bmatrix} c_{00} & c_{01} \\ c_{10} & c_{11} \end{bmatrix} = \begin{bmatrix} a_{00} & a_{01} \\ a_{10} & a_{11} \end{bmatrix} * \begin{bmatrix} b_{00} & b_{01} \\ b_{10} & b_{11} \end{bmatrix}$$
$$= \begin{bmatrix} m_1 + m_4 - m_5 + m_7 & m_3 + m_5 \\ m_2 + m_4 & m_1 + m_3 - m_2 + m_6 \end{bmatrix},$$

$$m_1 = (a_{00} + a_{11}) * (b_{00} + b_{11}),$$

$$m_2 = (a_{10} + a_{11}) * b_{00},$$

$$m_3 = a_{00} * (b_{01} - b_{11}),$$

$$m_4 = a_{11} * (b_{10} - b_{00}),$$

$$m_5 = (a_{00} + a_{01}) * b_{11},$$

$$m_6 = (a_{10} - a_{00}) * (b_{00} + b_{01}),$$

$$m_7 = (a_{01} - a_{11}) * (b_{10} + b_{11}),$$

7 Multiplications & 18 additions

Strassen's Matrix Multiplication - Analysis

Since $n = 2^k$,

$$\begin{aligned} M(2^k) &= 7M(2^{k-1}) = 7[7M(2^{k-2})] = 7^2M(2^{k-2}) = \dots \\ &= 7^i M(2^{k-i}) \dots = 7^k M(2^{k-k}) = 7^k. \end{aligned}$$

Since $k = \log_2 n$,

$$M(n) = 7^{\log_2 n} = n^{\log_2 7} \approx n^{2.807}.$$