



Unit III – Dynamic Programming

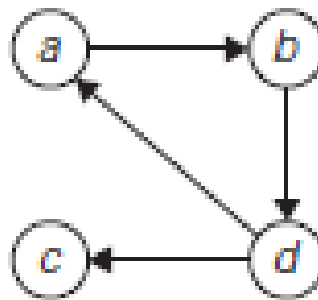
- Dynamic Programming
 - Computing a Binomial Coefficient
 - **Warshall's algorithm**
 - Floyd's algorithm
 - Optimal Binary Search Trees
 - Knapsack Problem and Memory functions



Warshall's algorithm



- Compute the Transitive closure of a directed graph
- The **transitive closure** of a directed graph with n vertices can be defined as the $n \times n$ boolean matrix $T = \{t_{ij}\}$, in which the element in the i th row and the j th column is 1 if there exists a nontrivial path (i.e., directed path of a positive length) from the i th vertex to the j th vertex; otherwise, t_{ij} is 0.
- *Example: directed graph (digraph)*

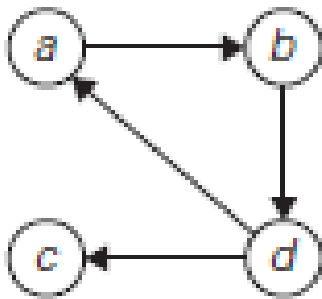




Warshall's algorithm



- **Adjacency Matrix** - $A = \{a_{ij}\}$ of a digraph is the boolean matrix that has 1 in the i^{th} row and j^{th} column if and only if there is a directed edge from i^{th} vertex to j^{th} vertex.



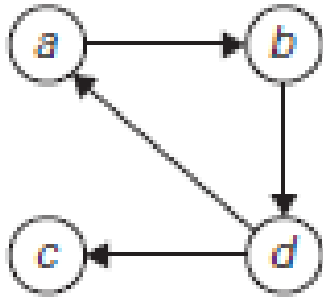
$$A = \begin{matrix} & \begin{matrix} a & b & c & d \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \end{matrix} & \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 \end{bmatrix} \end{matrix}$$



Warshall's algorithm



- Transitive closure



	a	b	c	d
a	1	1	1	1
b	1	1	1	1
c	0	0	0	0
d	1	1	1	1



Warshall's algorithm

- series of $n \times n$ boolean matrices: $R(0), \dots, R(k-1), R(k), \dots R(n)$.
- Matrix value is 1 – using the formula as follows

$$r_{ij}^{(k)} = r_{ij}^{(k-1)} \text{ or } \left(r_{ik}^{(k-1)} \text{ and } r_{kj}^{(k-1)} \right)$$

$$R^{(k-1)} = \begin{matrix} & \begin{matrix} j & k \end{matrix} \\ \begin{matrix} i \\ \uparrow \\ 0 \end{matrix} \rightarrow & \begin{bmatrix} & & \\ & 1 & \\ & & \end{bmatrix} \end{matrix} \implies R^{(k)} = \begin{matrix} & \begin{matrix} j & k \end{matrix} \\ \begin{matrix} i \\ \uparrow \\ 0 \end{matrix} \rightarrow & \begin{bmatrix} & & \\ & 1 & \\ & 1 & 1 \end{bmatrix} \end{matrix}$$

Rule for changing zeros in Warshall's algorithm.

Warshall's algorithm – Example

$$A = \begin{matrix} & \begin{matrix} a & b & c & d \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \end{matrix} & \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 \end{bmatrix} \end{matrix}$$

R_0	a	b	c	D
a	0	1	0	0
b	0	0	0	1
c	0	0	0	0
d	1	0	1	0

R_1	a	b	c	D
a	0	1	0	0
b	0	0	0	1
c	0	0	0	0
d	1	1	1	0

R_2	a	b	c	D
a	0	1	0	1
b	0	0	0	1
c	0	0	0	0
d	1	1	1	1

R_3	a	b	c	D
a	0	1	0	1
b	0	0	0	1
c	0	0	0	0
d	1	1	1	1

R_4	a	b	c	D
a	1	1	1	1
b	1	1	1	1
c	0	0	0	0
d	1	1	1	1



Warshall's algorithm - Algorithm

ALGORITHM *Warshall*($A[1..n, 1..n]$)

//Implements Warshall's algorithm for computing the transitive closure|

//Input: The adjacency matrix A of a digraph with n vertices

//Output: The transitive closure of the digraph

$R^{(0)} \leftarrow A$

for $k \leftarrow 1$ **to** n **do**

for $i \leftarrow 1$ **to** n **do**

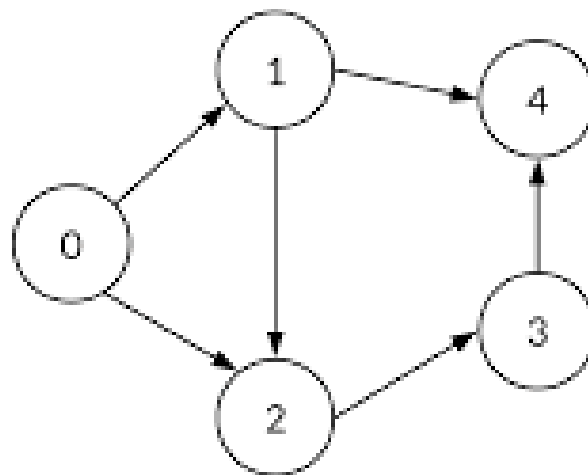
for $j \leftarrow 1$ **to** n **do**

$R^{(k)}[i, j] \leftarrow R^{(k-1)}[i, j] \text{ or } (R^{(k-1)}[i, k] \text{ and } R^{(k-1)}[k, j])$

return $R^{(n)}$



Warshall's algorithm - Example



Adjacency Matrix

	0	1	2	3	4
0	0	1	1	0	0
1	0	0	1	0	1
2	0	0	0	1	0
3	0	0	0	0	1
4	0	0	0	0	0