



23CST202 – Operating Systems

Paging Structure of Page Table

Dr.V.Savitha Associate Professor SNS College of Technology







- Hierarchical Paging
- Hashed Page Tables
- Inverted Page Tables



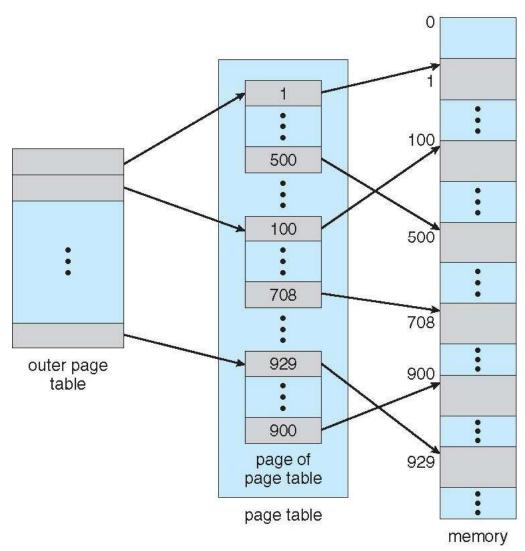


- Break up the logical address space into multiple page tables
- A simple technique is a two-level page table



Two-Level Page-Table Scheme





23CST202- OS-Structure of Page Table-V.Savitha, AP/CSE



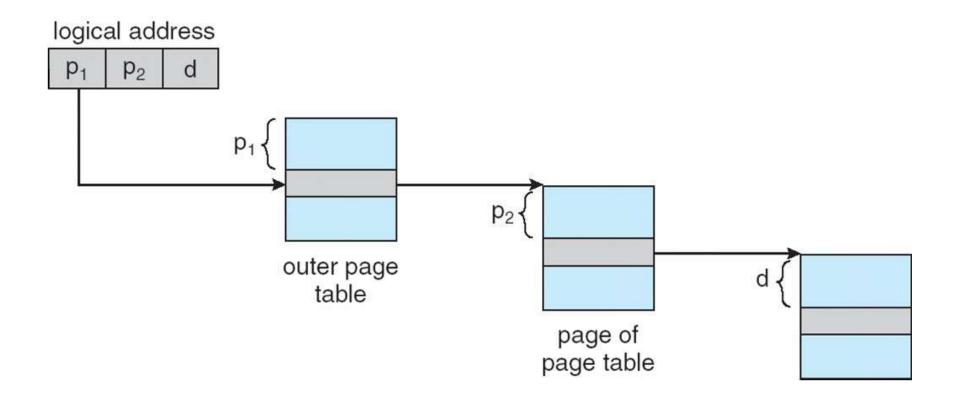


- A logical address (on 32-bit machine with 1K page size) is divided into:
 - a page number consisting of 22 bits
 - a page offset consisting of 10 bits
- A logical address (on 32-bit machine with 1K page size) is divided into:
 - a page number consisting of 22 bits
 - a page offset consisting of 10 bits

page number		page offset	
<i>p</i> ₁	<i>p</i> ₂	d	
12	10	10	

*p*₁ is an index into the outer page table, and *p*₂ is the displacement within the page of the inner page table known as forward-mapped page table









- Two-level paging scheme not sufficient
- Outer page table has 2⁴² entries or 2⁴⁴ bytes

outer page	inner page	page offset	
P ₁	<i>p</i> ₂	d	
42	10	12	

• Three-level Paging Scheme

2nd outer page	outer page	inner page	offset
p_1	p_2	p_3	d
32	10	10	12





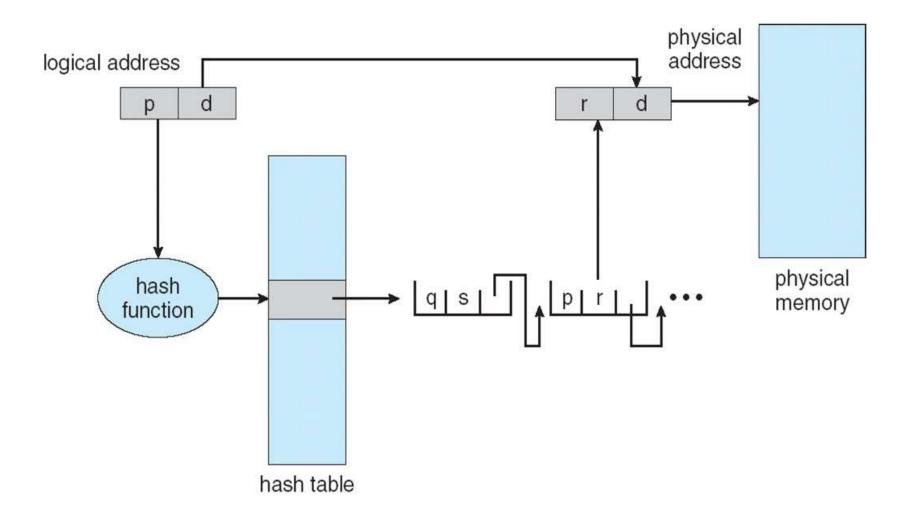


- Virtual page number is **hashed** into a page table
- Each element contains
 (1) the virtual page number
 (2) the value of the mapped page frame
 (3) a pointer to the next element
- Virtual page numbers are compared in this chain searching for a match
 - If a match is found, the corresponding physical frame is extracted



Hashed Page Table







Inverted Page Table



 Rather than each process having a page table and keeping track of all possible logical pages, track all physical pages

