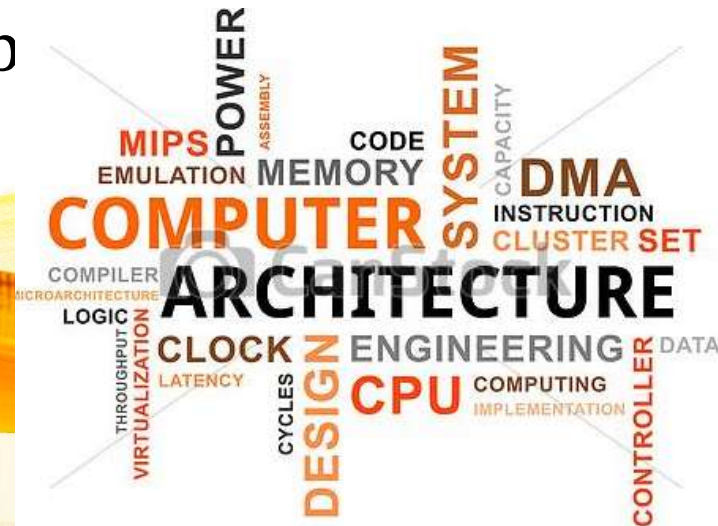


# UNIT II

## ARITHMETIC OPERATIONS

Addition and subtraction of signed numbers – **Design of fast adders** –  
Multiplication of positive numbers - Signed operand multiplication- fast  
multiplication – Integer division – Floating point numb



# Recap the previous Class



# Binary Adders

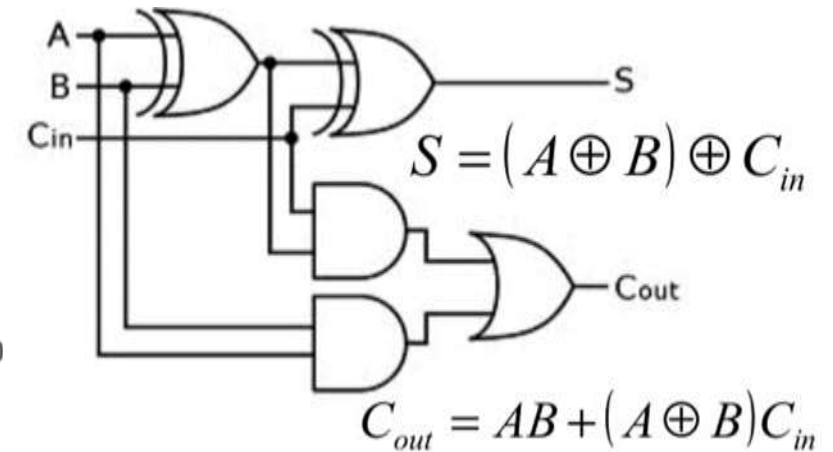
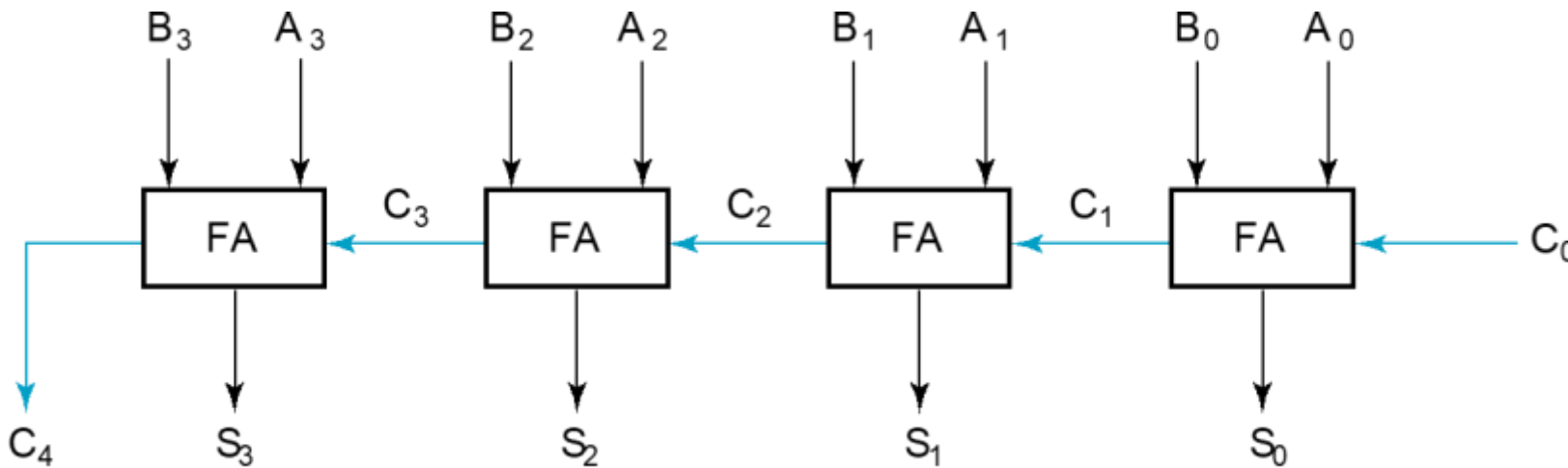
## Note:

carry out of cell  $i$  becomes carry in of cell  $i + 1$

Description	Subscript 3 2 1 0	Name
Carry In	0 1 1 0	$C_i$
Augend	1 0 1 1	$A_i$
Addend	<u>0 0 1 1</u>	$B_i$
Sum	1 1 1 0	$S_i$
Carry out	0 0 1 1	$C_{i+1}$

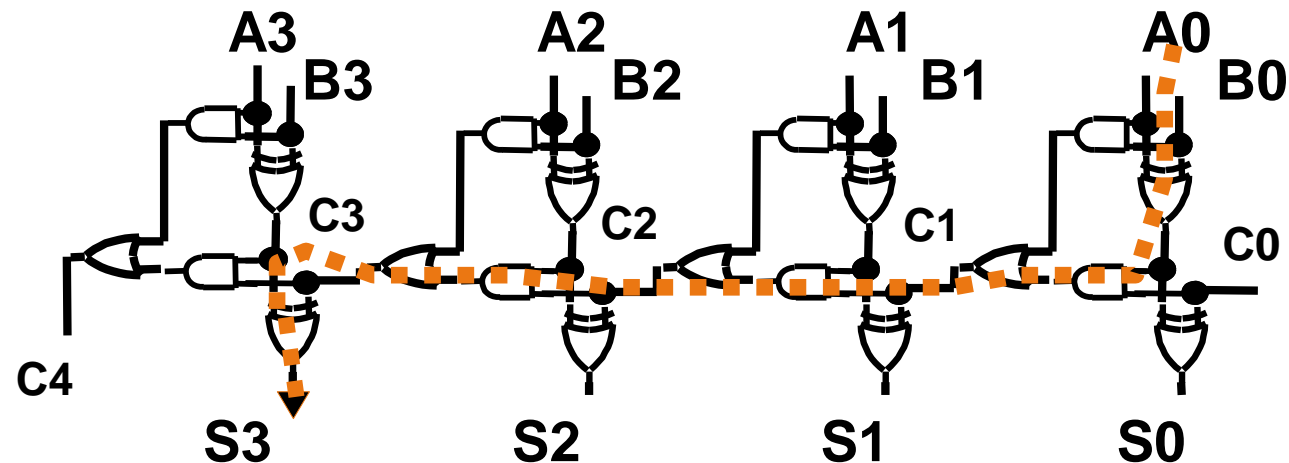
# 4 bit Ripple carry Adder

- A four-bit Ripple Carry Adder made from four 1-bit Full Adders



# Carry Propagation & Delay

- One problem with the addition of binary numbers is the length of time to propagate the ripple carry from the least significant bit to the most significant bit.
- The gate-level propagation path for a 4-bit ripple carry adder





# Carry Lookahead Adder

$$S_i = x_i \oplus y_i \oplus c_i$$

$$C_{i+1} = x_i y_i + x_i c_i + y_i c_i$$

Factorizing

$$C_{i+1} = x_i y_i + (x_i + y_i) c_i$$

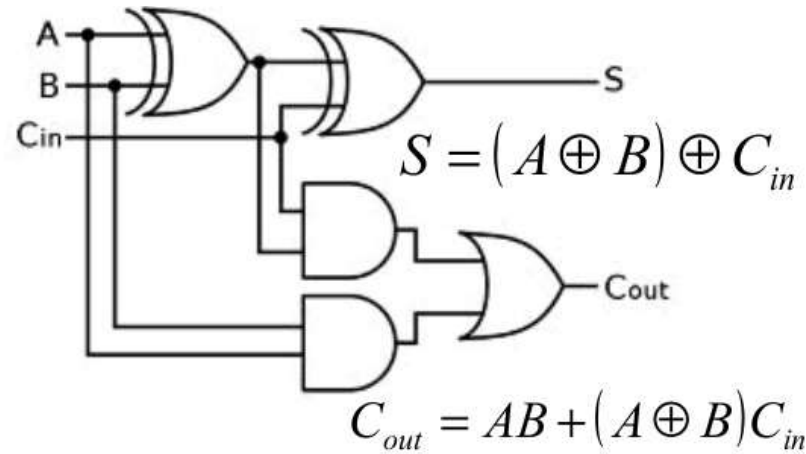
We can write

$$C_{i+1} = G_i + P_i c_i$$

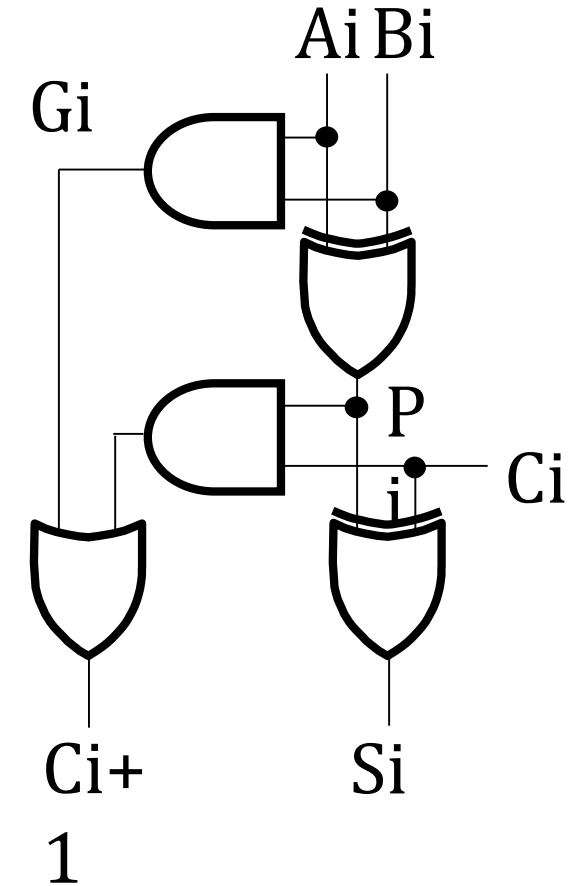
Where

$$G_i = x_i y_i$$

$$P_i = x_i + y_i$$



These two signal conditions are called **generate**, denoted as  **$G_i$** , and **propagate**, denoted as  **$P_i$**  respectively



# Carry Lookahead Adder

- In the ripple carry adder:
  - $G_i$ ,  $P_i$ , and  $S_i$  are local to each cell of the adder
  - $C_i$  is also local each cell
- In the carry lookahead adder, in order to reduce the length of the carry chain,  $C_i$  is changed to a more global function spanning multiple cells
- Defining the equations for the Full Adder in term of the  $P_i$  and  $G_i$ :

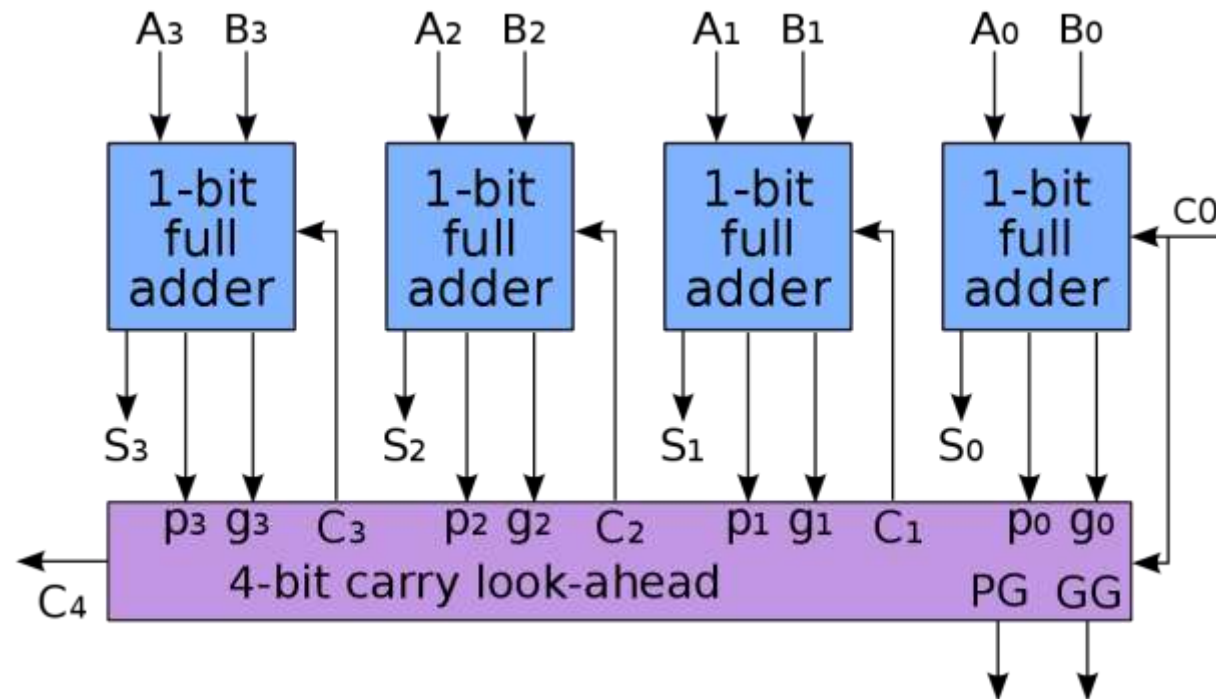
$$P_i = A_i \oplus B_i$$

$$G_i = A_i B_i$$

$$S_i = P_i \oplus C_i$$

$$C_{i+1} = G_i + P_i C_i$$

# Carry Lookahead Adder



$$C_1 = G_0 + P_0 C_0$$

$$C_2 = G_1 + P_1 C_1 = G_1 + P_1(G_0 + P_0 C_0) \\ = G_1 + P_1 G_0 + P_1 P_0 C_0$$

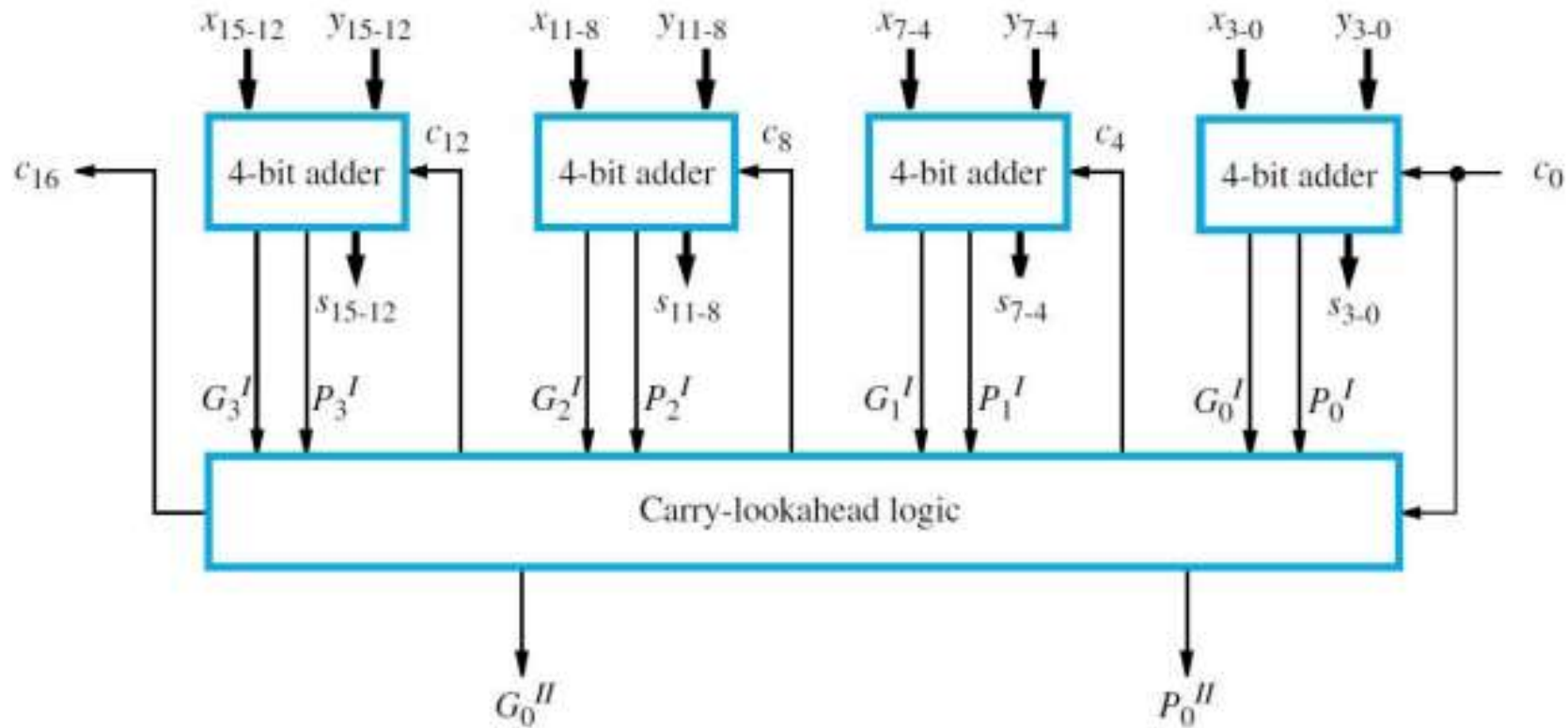
$$C_3 = G_2 + P_2 C_2 = G_2 + P_2(G_1 + P_1 G_0 + P_1 P_0 C_0) \\ = G_2 + P_2 G_1 + P_2 P_1 G_0 + P_2 P_1 P_0 C_0$$

$$C_4 = G_3 + P_3 C_3 = G_3 + P_3 G_2 + P_3 P_2 G_1 \\ + P_3 P_2 P_1 G_0 + P_3 P_2 P_1 P_0 C_0$$

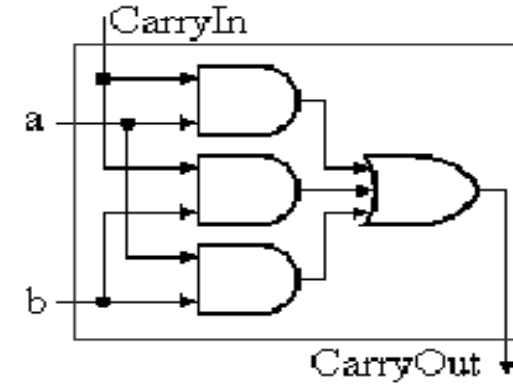
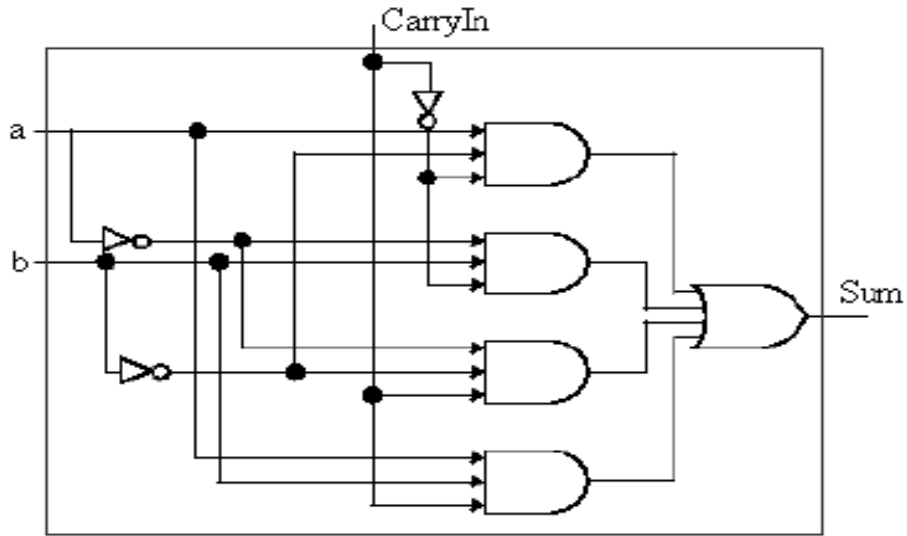
$$C_{i+1} = G_i + P_i G_{i-1} + P_i P_{i-1} C_{i-1}$$



# 16 bit Carry Lookahead Adder



# Assessment



$$\text{Carryout} = (b \cdot \text{CarryIn}) + (a \cdot \text{CarryIn}) + (a \cdot b)$$

$$\text{Sum} = (a \cdot b' \cdot \text{CarryIn}') + (a' \cdot b \cdot \text{CarryIn}') + (a' \cdot b' \cdot \text{CarryIn}) + (a \cdot b \cdot \text{CarryIn})$$



**sns**  
INSTITUTIONS



*Thank You*