

SNS COLLEGE OF TECHNOLOGY

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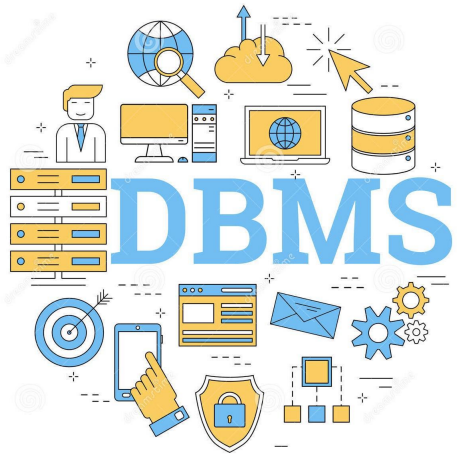
Department of Information Technology

19CST202 – Database Management System
II B.Tech. AIML/ IV SEMESTER

UNIT II : RELATIONAL MODEL

Topic 1 : Relational Data Model - keys, referential integrity and foreign keys

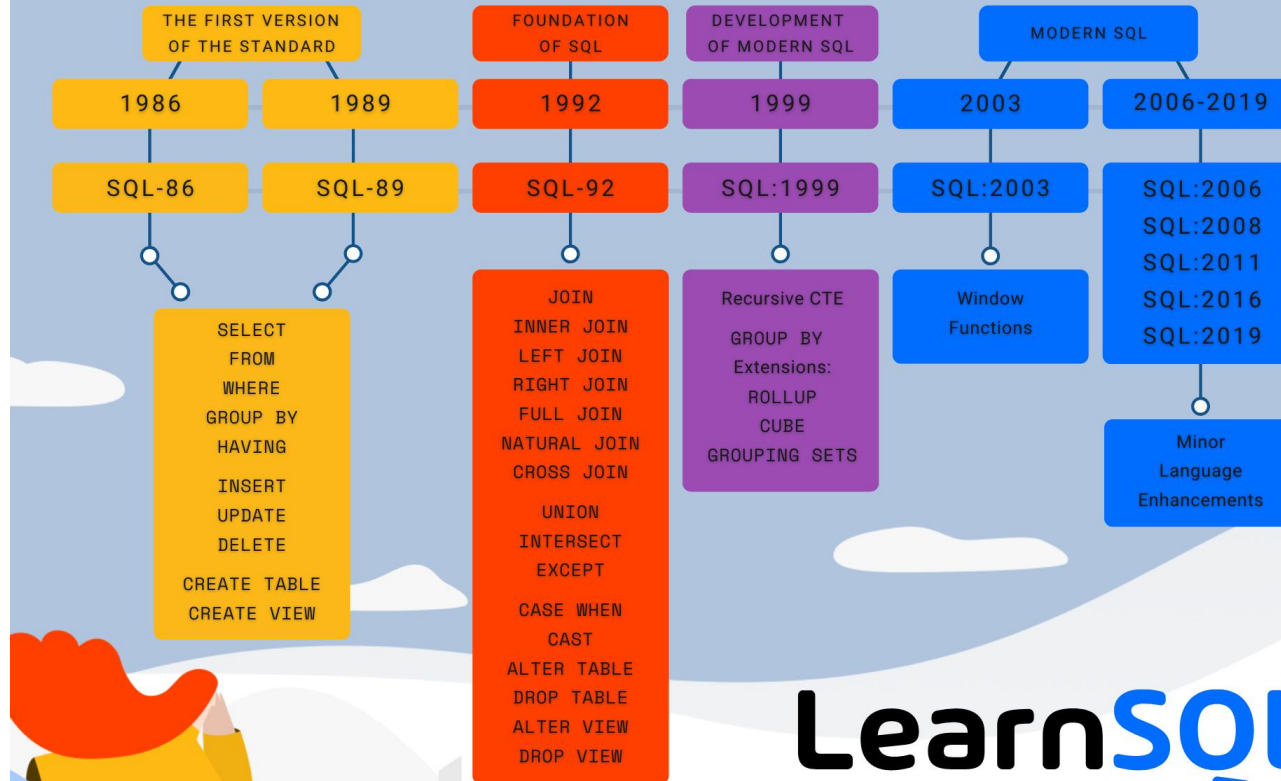
Relational Data Model - keys, referential integrity and foreign keys, Relational Algebra - SQL fundamentals- Introduction, data definition in SQL, table, key and foreign key definitions, update behaviors-Intermediate SQL-Advanced SQL features -Embedded SQL- Dynamic SQL, CASE Studies- Oracle: Database Design and Querying Tools; SQL Variations and Extensions





History

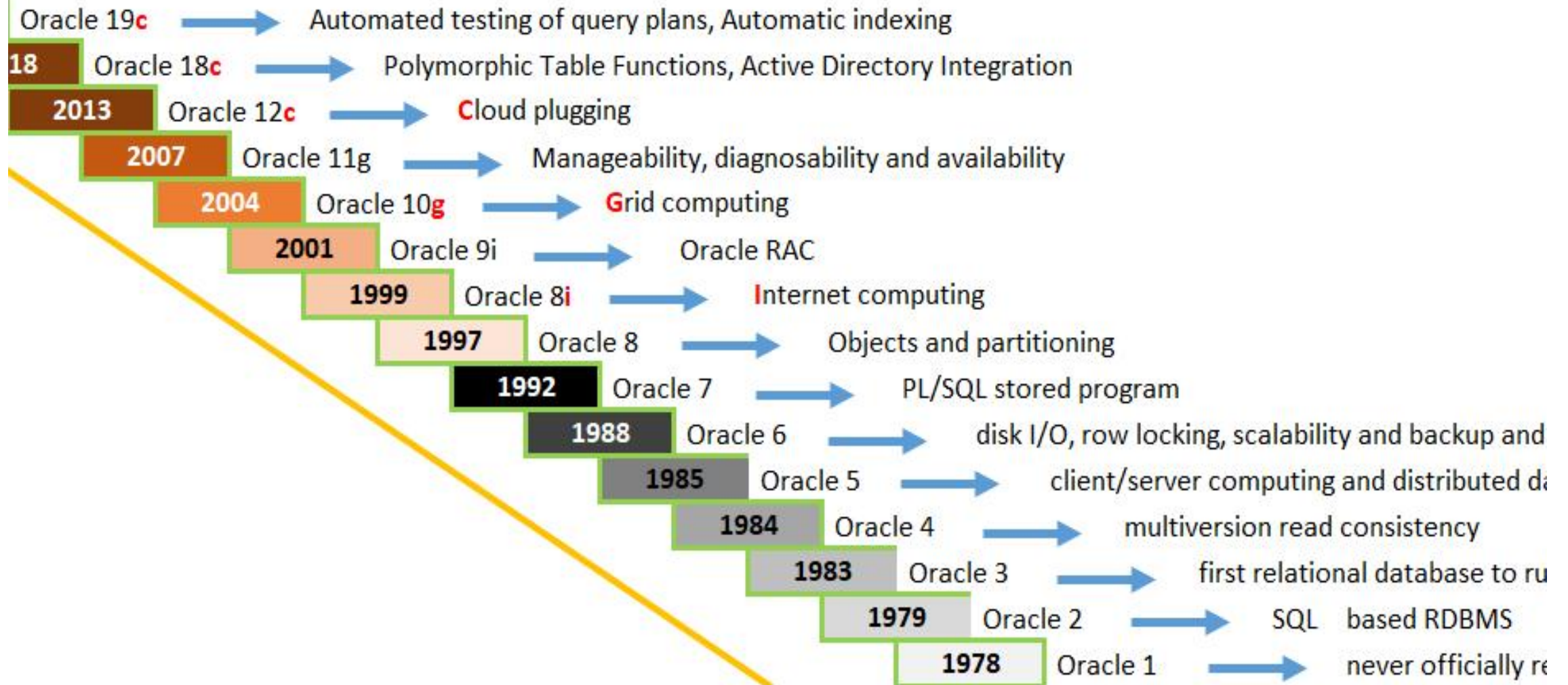
The History of SQL Standards



LearnSQL.com

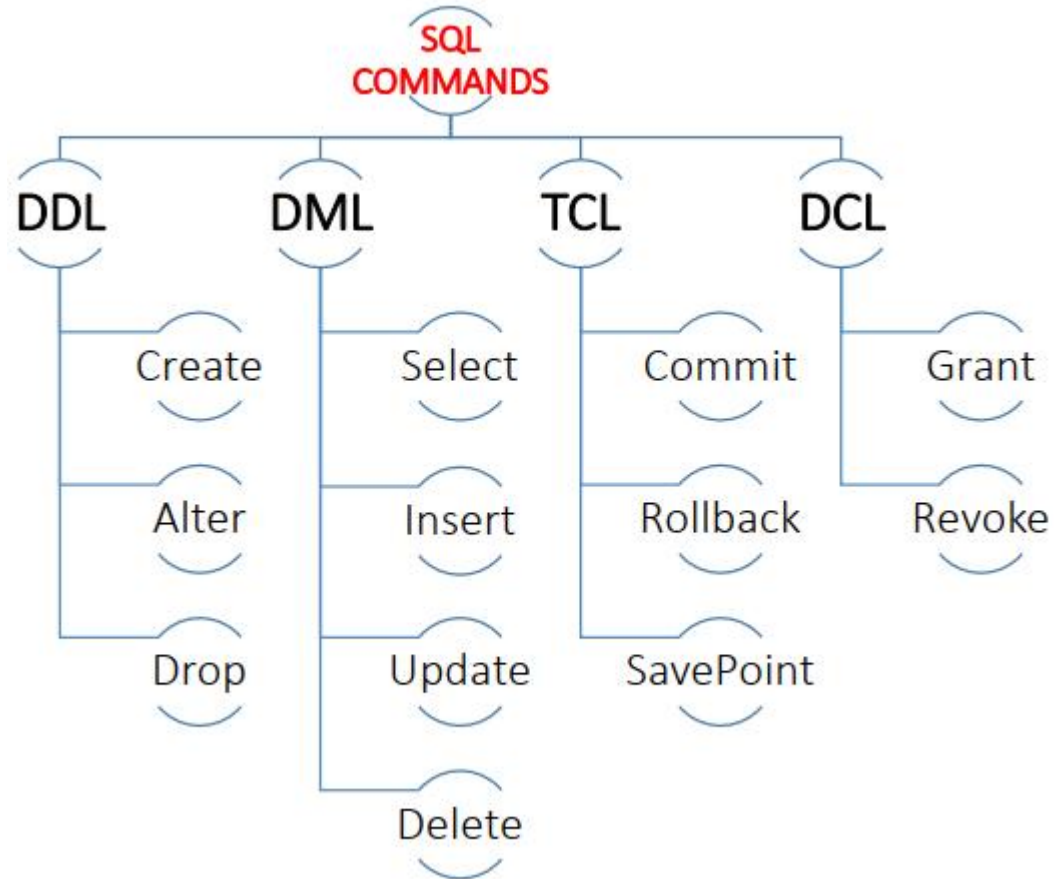


History of Oracle Database Versions





SQL Parts



- DML - Data Manipulation Language
- DDL - Data Definition Language
- DCL - Data Control Language
- TCL - Transaction Control Language
- DQL - Data Query Language

Data Definition Language

- The SQL data-definition language (DDL) allows the specification of information about relations, including:
 - The schema for each relation.
 - The type of values associated with each attribute.
 - The Integrity constraints
 - The set of indices to be maintained for each relation.
 - Security and authorization information for each relation.
 - The physical storage structure of each relation on disk.

- **char(*n*)**. Fixed length character string, with user-specified length *n*.
- **varchar(*n*)**. Variable length character strings, with user-specified maximum length *n*.
- **int**. Integer (a finite subset of the integers that is machine-dependent).
- **smallint**. Small integer (a machine-dependent subset of the integer domain type).
- **numeric(*p*,*d*)**. Fixed point number, with user-specified precision of *p* digits, with *d* digits to the right of decimal point. (ex., **numeric(3,1)**, allows 44.5 to be stores exactly, but not 444.5 or 0.32)
- **real, double precision**. Floating point and double-precision floating point numbers, with machine-dependent precision.
- **float(*n*)**. Floating point number, with user-specified precision of at least *n* digits.

- An SQL relation is defined using the **create table** command:

create table r

$(A_1 D_1, A_2 D_2, \dots, A_n D_n$
 $(\text{integrity-constraint}_1),$
 $\dots,$
 $(\text{integrity-constraint}_k))$

```
create table instructor (  
  ID           char(5),  
  name        varchar(20),  
  dept_name   varchar(20),  
  salary      numeric(8,2))
```

- r is the name of the relation
- each A_i is an attribute name in the schema of relation r
- D_i is the data type of values in the domain of attribute A_i



Integrity Constraints in Create Table

- Types of integrity constraints
 - **primary key** (A_1, \dots, A_n)
 - **foreign key** (A_m, \dots, A_n) **references** r
 - **not null**
- SQL prevents any update to the database that violates an integrity constraint.



Integrity Constraints in Create Table

```
create table instructor (  
    ID          char(5),  
    name       varchar(20) not null,  
    dept_name  varchar(20),  
    salary     numeric(8,2),  
    primary key (ID),  
    foreign key (dept_name) references department);
```



Few More Relation Definitions

create table *student* (

ID **varchar(5),**

name **varchar(20) not null,**

dept_name **varchar(20),**

tot_cred **numeric(3,0),**

primary key (*ID*),

foreign key (*dept_name*) references *department*);



Few More Relation Definitions

- **create table** *takes* (

ID **varchar(5),**

course_id **varchar(8),**

sec_id **varchar(8),**

semester **varchar(6),**

year **numeric(4,0),**

grade **varchar(2),**

primary key (*ID, course_id, sec_id, semester, year*) ,

foreign key (*ID*) **references** *student*,

foreign key (*course_id, sec_id, semester, year*) **references** *section*);

Few More Relation Definitions

- **create table** *course* (

 course_id **varchar**(8),

 title **varchar**(50),

 dept_name **varchar**(20),

 credits **numeric**(2,0),

 primary key (*course_id*),

 foreign key (*dept_name*) **references** *department*);

- **Insert**
 - **insert into** *instructor* **values** ('10211', 'Smith', 'Biology', 66000);
- **Delete**
 - Remove all tuples from the *student* relation
 - **delete from** *student*
- **Drop Table**
 - **drop table** *r*
- **Alter**
 - **alter table** *r* **add** *A D*
 - where *A* is the name of the attribute to be added to relation *r* and *D* is the domain of *A*.
 - All exiting tuples in the relation are assigned *null* as the value for the new attribute.
 - **alter table** *r* **drop** *A*
 - where *A* is the name of an attribute of relation *r*
 - Dropping of attributes not supported by many databases.

Basic Query Structure

- A typical SQL query has the form:

select A_1, A_2, \dots, A_n

from r_1, r_2, \dots, r_m

where P

A_i represents an attribute

- R_i represents a relation
 - P is a predicate.
- The result of an SQL query is a relation.

- The **select** clause lists the attributes desired in the result of a query
 - corresponds to the projection operation of the relational algebra

- **Example: find the names of all instructors:**

select *name*

from *instructor*

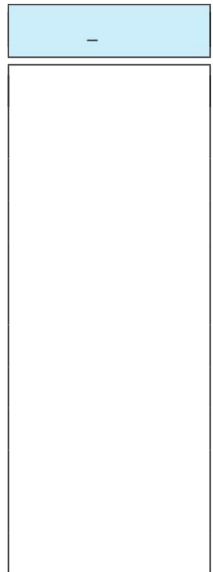
- NOTE: SQL names are case insensitive (i.e., you may use upper- or lower-case letters.)
 - E.g., *Name* \equiv *NAME* \equiv *name*
 - Some people use upper case wherever we use bold font.

- **SQL allows duplicates in relations as well as in query results.**
- To force the elimination of duplicates, insert the keyword **distinct** after select.
- Find the department names of all instructors, and remove duplicates

```
select distinct dept_name  
from instructor
```

- The keyword **all** specifies that duplicates should not be removed.

```
select all dept_name  
from instructor
```



- An asterisk in the select clause denotes “all attributes”

s e l e c t *

from *instructor*

- An attribute can be a literal with no **from** clause

select '437'

- Results is a table with one column and a single row with value “437”
- Can give the column a name using:

select '437' *as* *FOO*

- An attribute can be a literal with **from** clause

select 'A' **from** *instructor*

- Result is a table with one column and N rows (number of tuples in the *instructors* table), each row with value “A”

- The **select** clause can contain arithmetic expressions involving the operation, +, -, *, and /, and operating on constants or attributes of tuples.
- The query:

```
select ID, name, salary/12  
from instructor
```

would return a relation that is the same as the *instructor* relation, except that the value of the attribute *salary* is divided by 12.

- Can rename “*salary/12*” using the **as** clause:
select ID, name, *salary/12 as monthly_salary*

- The **where** clause specifies conditions that the result must satisfy
 - Corresponds to the selection predicate of the relational algebra.
- To find all instructors in Comp. Sci. dept

```
select name  
from instructor  
where dept_name = 'Comp. Sci.'
```

- SQL allows the use of the logical connectives **and, or, and not**
- The operands of the logical connectives can be expressions involving the comparison operators **<, <=, >, >=, =, and <>**.
- Comparisons can be applied to results of arithmetic expressions
- To find all instructors in Comp. Sci. dept with salary > 70000

```
select name  
from instructor  
where dept_name = 'Comp. Sci.' and salary > 70000
```



- The **from** clause lists the relations involved in the query
 - Corresponds to the Cartesian product operation of the relational algebra.
- Find the Cartesian product *instructor X teaches*

select *

from *instructor, teaches*

- generates every possible instructor – teaches pair, with all attributes from both relations.
- For common attributes (e.g., *ID*), the attributes in the resulting table are renamed using the relation name (e.g., *instructor.ID*)
- Cartesian product not very useful directly, but useful combined with where-clause condition (selection operation in relational algebra).

- Find the names of all instructors who have taught some course and the course_id
 - **select** *name, course_id*
from *instructor, teaches*
where *instructor.ID = teaches.ID*
- Find the names of all instructors in the Art department who have taught some course and the course_id
 - **select** *name, course_id*
from *instructor, teaches*
where *instructor.ID = teaches.ID*
and *instructor.dept_name = 'Art'*

	-

The Rename Operation

- The SQL allows renaming relations and attributes using the **as** clause:

old-name as new-name

- Find the names of all instructors who have a higher salary than some instructor in 'Comp. Sci'.

- **select distinct** *T.name*
from *instructor as T, instructor as S*
where *T.salary > S.salary and S.dept_name = 'Comp. Sci.'*

- Keyword **as** is optional and may be omitted
instructor as T \equiv *instructor T*

- SQL includes a string-matching operator for comparisons on character strings. The operator **like** uses patterns that are described using two special characters:
 - percent (%). The % character matches any substring.
 - underscore (_). The _ character matches any character.
- Find the names of all instructors whose name includes the substring “dar”.

```
select name  
from instructor  
where name like '%dar%'
```

- Match the string “100%”

```
like '100 \% ' escape '\'
```

in that above we use backslash (\) as the escape character.

<https://www.geeksforgeeks.org/sql-string-functions/>

- Patterns are case sensitive.
- Pattern matching examples:
 - 'Intro%' matches any string beginning with “Intro”.
 - '%Comp%' matches any string containing “Comp” as a substring.
 - '___' matches any string of exactly three characters.
 - '___ %' matches any string of at least three characters.
- SQL supports a variety of string operations such as
 - concatenation (using “||”)
 - converting from upper to lower case (and vice versa)
 - finding string length, extracting substrings, etc.



- List in alphabetic order the names of all instructors

select distinct *name*

from *instructor*

order by *name*

- We may specify **desc** for descending order or **asc** for ascending order, for each attribute; ascending order is the default.
 - Example: **order by** *name* **desc**
- Can sort on multiple attributes
 - Example: **order by** *dept_name, name*

- SQL includes a **between** comparison operator
- Example: Find the names of all instructors with salary between \$90,000 and \$100,000 (that is, \geq \$90,000 and \leq \$100,000)
 - **select** *name*
from *instructor*
where *salary* **between** 90000 **and** 100000
- Tuple comparison
 - **select** *name, course_id*
from *instructor, teaches*
where (*instructor.ID, dept_name*) = (*teaches.ID, 'Biology'*);

- Find courses that ran in Fall 2017 or in Spring 2018

(select *course_id* from *section* where *sem* = 'Fall' and *year* = 2017)

union

(select *course_id* from *section* where *sem* = 'Spring' and *year* = 2018)

- Find courses that ran in Fall 2017 and in Spring 2018

(select *course_id* from *section* where *sem* = 'Fall' and *year* = 2017)

intersect

(select *course_id* from *section* where *sem* = 'Spring' and *year* = 2018)

Find courses that ran in Fall 2017 but not in Spring 2018

(select *course_id* from *section* where *sem* = 'Fall' and *year* = 2017)

except

(select *course_id* from *section* where *sem* = 'Spring' and *year* = 2018)

Set Operations (Cont.)

- Set operations **union**, **intersect**, and **except**
 - Each of the above operations automatically eliminates duplicates
- To retain all duplicates use the
 - **union all**,
 - **intersect all**
 - **except all**.



Null Values

- It is possible for tuples to have a null value, denoted by **null**, for some of their attributes
- **null** signifies an unknown value or that a value does not exist.
- The result of any arithmetic expression involving **null** is **null**
 - Example: $5 + \mathbf{null}$ returns **null**
- The predicate **is null** can be used to check for null values.
 - Example: Find all instructors whose salary is null.

select *name*

from *instructor*

where *salary is null*

- The predicate **is not null** succeeds if the value on which it is applied is not null.

Aggregate Functions

- These functions operate on the multiset of values of a column of a relation, and return a value

avg: average value

min: minimum value

max: maximum value

sum: sum of values

count: number of values



- Find the average salary of instructors in the Computer Science department
 - **select avg** (*salary*)
from *instructor*
where *dept_name*= 'Comp. Sci.';
- Find the total number of instructors who teach a course in the Spring 2018 semester
 - **select count** (**distinct** *ID*)
from *teaches*
where *semester* = 'Spring' **and** *year* = 2018;
- Find the number of tuples in the *course* relation
 - **select count** (*)
from *course*;



- Find the average salary of instructors in each department
 - **select** *dept_name*, **avg** (*salary*) **as** *avg_salary*
from *instructor*
group by *dept_name*;

-	-

		-	

- Find the names and average salaries of all departments whose average salary is greater than 42000

```
select dept_name, avg (salary) as avg_salary
from instructor
group by dept_name
having avg (salary) > 42000;
```

- Note: predicates in the **having** clause are applied after the formation of groups whereas predicates in the **where** clause are applied before forming groups

- SQL provides a mechanism for the nesting of subqueries. A **subquery** is a **select-from-where** expression that is nested within another query.
- The nesting can be done in the following SQL query

```
select  $A_1, A_2, \dots, A_n$   
from  $r_1, r_2, \dots, r_m$   
where  $P$ 
```

as follows:

- **From clause:** r_i can be replaced by any valid subquery
- **Where clause:** P can be replaced with an expression of the form:

B <operation> (subquery)

B is an attribute and <operation> to be defined later.

- **Select clause:**

A_i can be replaced by a subquery that generates a single value.



Set Membership

- Find courses offered in Fall 2017 and in Spring 2018

```
select distinct course_id
from section
where semester = 'Fall' and year= 2017 and
       course_id in (select course_id
                     from section
                     where semester = 'Spring' and year= 2018);
```

- Find courses offered in Fall 2017 but not in Spring 2018

```
select distinct course_id
from section
where semester = 'Fall' and year= 2017 and
       course_id not in (select course_id
                        from section
                        where semester = 'Spring' and year= 2018);
```

- Name all instructors whose name is neither “Mozart” nor Einstein”

```
select distinct name  
from instructor  
where name not in ('Mozart', 'Einstein')
```

- Find the total number of (distinct) students who have taken course sections taught by the instructor with *ID* 10101

```
select count (distinct ID)  
from takes  
where (course_id, sec_id, semester, year) in  
      (select course_id, sec_id, semester, year  
       from teaches  
       where teaches.ID= 10101);
```

- Note: Above query can be written in a much simpler manner.
The formulation above is simply to illustrate SQL features



Set Comparison - "some" Clause

- Find names of instructors with salary greater than that of some (at least one) instructor in the Biology department.

```
select distinct T.name  
from instructor as T, instructor as S  
where T.salary > S.salary and S.dept name = 'Biology';
```

- Same query using > **some** clause

```
select name  
from instructor  
where salary > some (select salary  
                  from instructor  
                  where dept name = 'Biology');
```

- $F \text{ <comp> some } r \Leftrightarrow \exists t \in r \text{ such that } (F \text{ <comp> } t)$

Where <comp> can be: <, ≤, >, =, ≠

$$(5 < \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline 6 \\ \hline \end{array}) = \text{true} \quad (\text{read: } 5 < \text{some tuple in the relation})$$

$$(5 < \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline \end{array}) = \text{false}$$

$$(5 = \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline \end{array}) = \text{true}$$

$$(5 \neq \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline \end{array}) = \text{true (since } 0 \neq 5)$$

$(= \text{ some}) \equiv \text{in}$

However, $(\neq \text{ some}) \not\equiv \text{not in}$



Set Comparison – “all” Clause

- Find the names of all instructors whose salary is greater than the salary of all instructors in the Biology department.

```
select name  
from instructor  
where salary > all (select salary  
                    from instructor  
                    where dept name = 'Biology');
```




- $F \text{ <comp> all } r \Leftrightarrow \forall t \in r (F \text{ <comp> } t)$ **Definition of “all” Clause**

$$(5 < \text{all } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline 6 \\ \hline \end{array}) = \text{false}$$

$$(5 < \text{all } \begin{array}{|c|} \hline 6 \\ \hline 10 \\ \hline \end{array}) = \text{true}$$

$$(5 = \text{all } \begin{array}{|c|} \hline 4 \\ \hline 5 \\ \hline \end{array}) = \text{false}$$

$$(5 \neq \text{all } \begin{array}{|c|} \hline 4 \\ \hline 6 \\ \hline \end{array}) = \text{true (since } 5 \neq 4 \text{ and } 5 \neq 6)$$

$(\neq \text{ all}) \equiv \text{not in}$
However, $(= \text{ all}) \not\equiv \text{in}$

Test for Empty Relations

- The **exists** construct returns the value **true** if the argument subquery is nonempty.
- **exists** $r \Leftrightarrow r \neq \emptyset$
- **not exists** $r \Leftrightarrow r = \emptyset$

Use of “exists” Clause

- Yet another way of specifying the query “Find all courses taught in both the Fall 2017 semester and in the Spring 2018 semester”

```
select course_id
from section as S
where semester = 'Fall' and year = 2017 and
      exists (select *
              from section as T
              where semester = 'Spring' and year = 2018
                  and S.course_id = T.course_id);
```

- **Correlation name** – variable S in the outer query
- **Correlated subquery** – the inner query



- Find all students who have taken all courses offered in the Biology department.

Use of "not exists" Clause

```
select distinct S.ID, S.name  
from student as S  
where not exists ( (select course_id  
                    from course  
                    where dept_name = 'Biology')  
except  
                (select T.course_id  
                 from takes as T  
                 where S.ID = T.ID));
```

- First nested query lists all courses offered in Biology
- Second nested query lists all courses a particular student took
- Note that $X - Y = \emptyset \Leftrightarrow X \subseteq Y$
- Note: Cannot write this query using = all and its variants



- The **unique** construct tests whether a subquery has any duplicate tuples in its result.
- The **unique** construct evaluates to “true” if a given subquery contains no duplicates .
- Find all courses that were offered at most once in 2017

```
select T.course_id  
from course as T  
where unique ( select R.course_id  
                from section as R  
                where T.course_id= R.course_id  
                    and R.year = 2017);
```



- SQL allows a subquery expression to be used in the **from** clause
- Find the average instructors' salaries of those departments where the average salary is greater than \$42,000."

```
select dept_name, avg_salary  
from ( select dept_name, avg (salary) as avg_salary  
      from instructor  
      group by dept_name)  
where avg_salary > 42000;
```

- Note that we do not need to use the **having** clause
- Another way to write above query

```
select dept_name, avg_salary  
from ( select dept_name, avg (salary)  
      from instructor  
      group by dept_name)  
      as dept_avg (dept_name, avg_salary)  
where avg_salary > 42000;
```




With Clause

- The **with** clause provides a way of defining a temporary relation whose definition is available only to the query in which the **with** clause occurs.
- Find all departments with the maximum budget

```
with max_budget (value) as  
    (select max(budget)  
     from department)  
select department.name  
from department, max_budget  
where department.budget = max_budget.value;
```



- Find all departments where the total salary is greater than the average of the total salary at all departments

Complex Queries using With Clause

```
with dept_total (dept_name, value) as  
    (select dept_name, sum(salary)  
     from instructor  
     group by dept_name),  
dept_total_avg(value) as  
    (select avg(value)  
     from dept_total)  
select dept_name  
from dept_total, dept_total_avg  
where dept_total.value > dept_total_avg.value;
```



- Scalar subquery is one which is used where a single value is expected
- List all departments along with the number of instructors in each department

```
select dept_name,  
        ( select count(*)  
          from instructor  
          where department.dept_name = instructor.dept_name)  
        as num_instructors  
from department;
```

- Runtime error if subquery returns more than one result tuple



- Deletion of tuples from a given relation.
- Insertion of new tuples into a given relation
- Updating of values in some tuples in a given relation



- Delete all instructors

delete from *instructor*

- Delete all instructors from the Finance department

delete from *instructor*

where *dept_name* = 'Finance';

- *Delete all tuples in the instructor relation for those instructors associated with a department located in the Watson building.*

delete from *instructor*

where *dept name* **in** (**select** *dept name*

from *department*

where *building* = 'Watson');

- Delete all instructors whose salary is less than the average salary of instructors

```
delete from instructor  
where salary < (select avg (salary)  
                from instructor);
```

- Problem: as we delete tuples from *instructor*, the average salary changes
- Solution used in SQL:
 1. First, compute **avg** (*salary*) and find all tuples to delete
 2. Next, delete all tuples found above (without recomputing **avg** or retesting the tuples)

- Add a new tuple to *course*

```
insert into course
```

```
values ('CS-437', 'Database Systems', 'Comp. Sci.', 4);
```

- or equivalently

```
insert into course (course_id, title, dept_name, credits)
```

```
values ('CS-437', 'Database Systems', 'Comp. Sci.', 4);
```

- Add a new tuple to *student* with *tot_creds* set to null

```
insert into student
```

```
values ('3003', 'Green', 'Finance', null);
```

- Make each student in the Music department who has earned more than 144 credit hours an instructor in the Music department with a salary of \$18,000.

insert into *instructor*

select *ID, name, dept_name, 18000*

from *student*

where *dept_name = 'Music' and total_cred > 144;*

- The **select from where** statement is evaluated fully before any of its results are inserted into the relation.

Otherwise queries like

insert into *table1 select * from table1*

would cause problem



- Give a 5% salary raise to all instructors

update instructor

set salary = salary * 1.05

- Give a 5% salary raise to those instructors who earn less than 70000

update instructor

set salary = salary * 1.05

where salary < 70000;

- Give a 5% salary raise to instructors whose salary is less than average

update instructor

set salary = salary * 1.05

**where salary < (select avg (salary)
from instructor);**

- Increase salaries of instructors whose salary is over \$100,000 by 3%, and all others by a 5%
 - Write two **update** statements:

```
update instructor  
  set salary = salary * 1.03  
  where salary > 100000;  
update instructor  
  set salary = salary * 1.05  
  where salary <= 100000;
```
 - The order is important
 - Can be done better using the **case** statement (next slide)

Case Statement for Conditional Updates

- Same query as before but with case statement

```
update instructor  
  set salary = case  
    when salary <= 100000 then salary * 1.05  
    else salary * 1.03  
  end
```



- Recompute and update `tot_creds` value for all students

```
update student S  
set tot_cred = (select sum(credits)  
from takes, course  
where takes.course_id = course.course_id and  
S.ID= takes.ID.and  
takes.grade <> 'F' and  
takes.grade is not null);
```

- Sets `tot_creds` to null for students who have not taken any course
- Instead of `sum(credits)`, use:

```
case  
when sum(credits) is not null then sum(credits)  
else 0  
end
```




TEXT BOOKS

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