



# **SNS COLLEGE OF TECHNOLOGY**

**Coimbatore-35.**

**An Autonomous Institution**

**Accredited by NBA – AICTE and Accredited by NAAC – UGC with ‘A++’ Grade  
Approved by AICTE, New Delhi & Affiliated to Anna University, Chennai**

**COURSE NAME : OPERATING SYSTEMS**

**II YEAR/ IV SEMESTER**

**UNIT – IV FILE SYSTEMS**

**Topic: Allocation Methods**

Dr.B.Vinodhini  
Associate Professor  
Department of Computer Science and Engineering



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## *Allocation Methods*

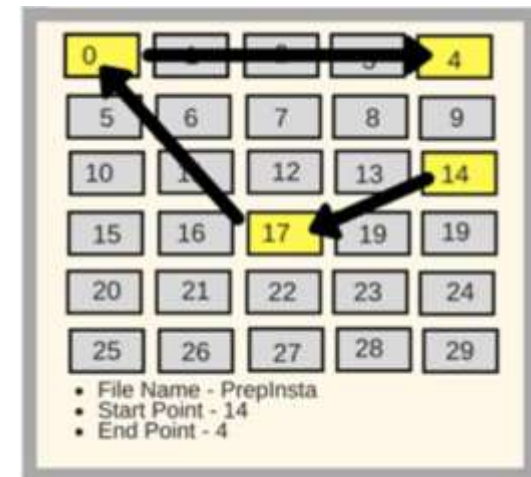


### *Allocation Methods*

*Contiguous Allocation*

*Linked Allocation*

*Indexed Allocation*





## *Allocation Methods*

- An allocation method refers to how disk blocks are allocated for files:
- Contiguous allocation
- Linked allocation
- Indexed allocation

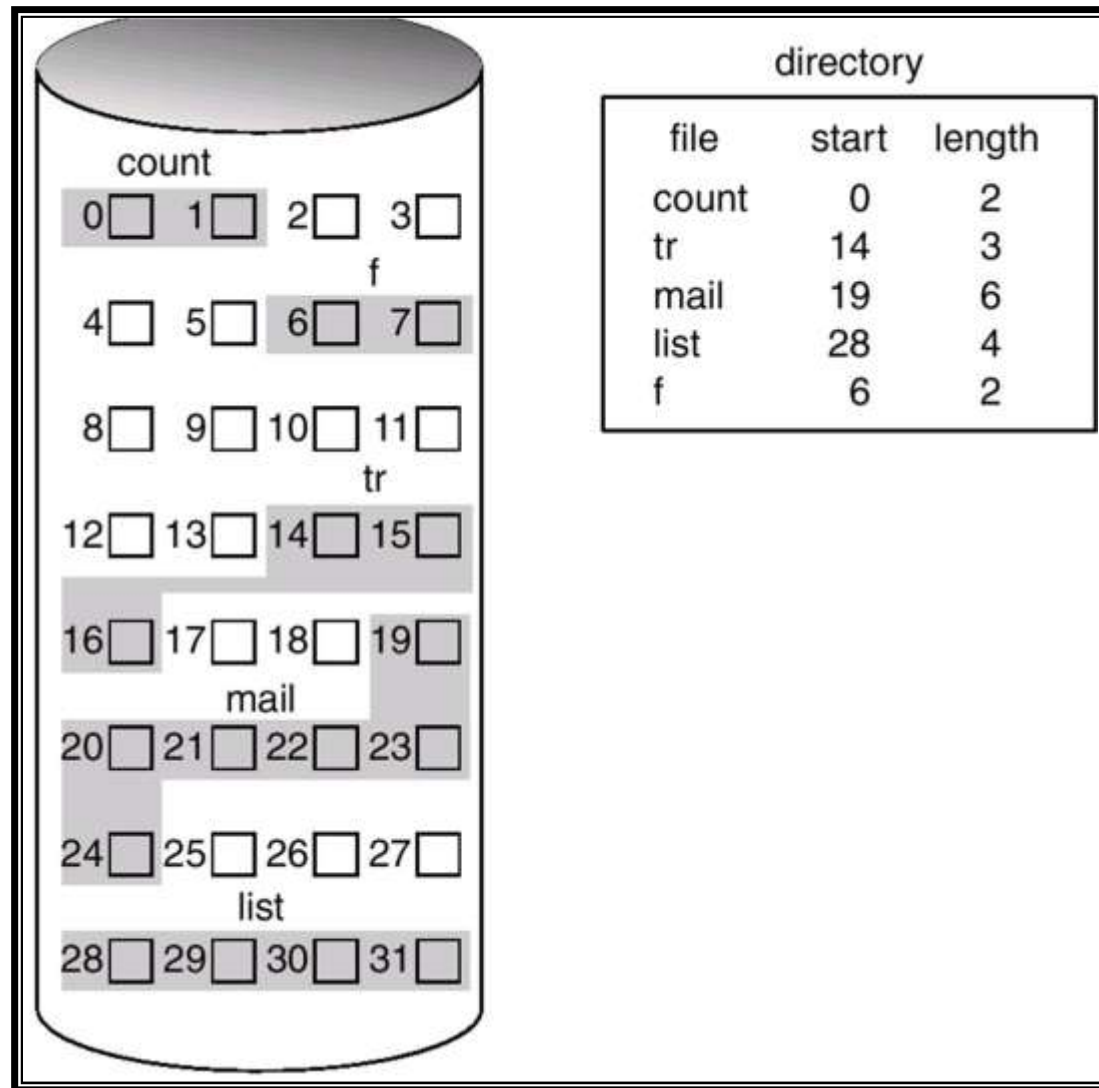


# *Contiguous Allocation*

- Each file occupies a set of contiguous blocks on the disk.
- Simple – only starting location (block #) and length (number of blocks) are required.
- Random access.
- Wasteful of space (dynamic storage-allocation problem).
- Files cannot grow.



# Contiguous Allocation of Disk Space





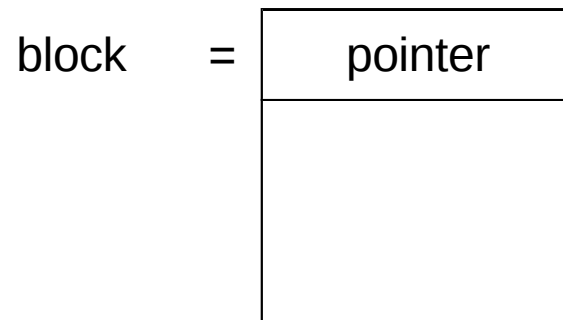
## *Extent-Based Systems*

- Many newer file systems (I.e. Veritas File System) use a modified contiguous allocation scheme.
- Extent-based file systems allocate disk blocks in **extents**.
- An **extent** is a contiguous block of disks. Extents are allocated for file allocation. A file consists of one or more extents.



## *Linked Allocation*

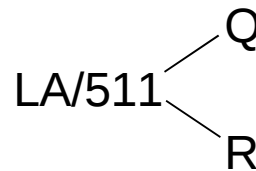
- Each file is a linked list of disk blocks: blocks may be scattered anywhere on the disk.





## *Linked Allocation (Cont.)*

- Simple – need only starting address
- Free-space management system – no waste of space
- No random access
- Mapping



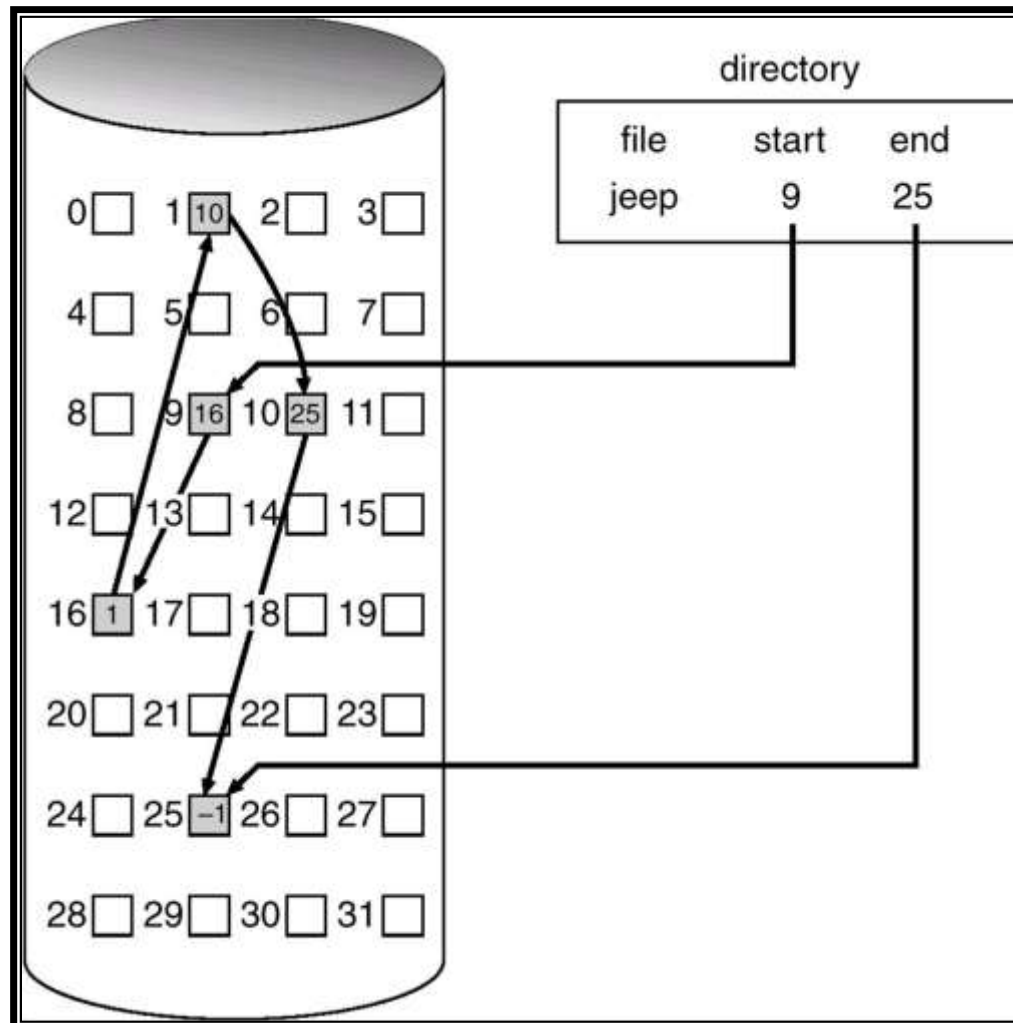
Block to be accessed is the Qth block in the linked chain of blocks representing the file.

Displacement into block =  $R + 1$

File-allocation table (FAT) – disk-space allocation used by MS-DOS and OS/2.

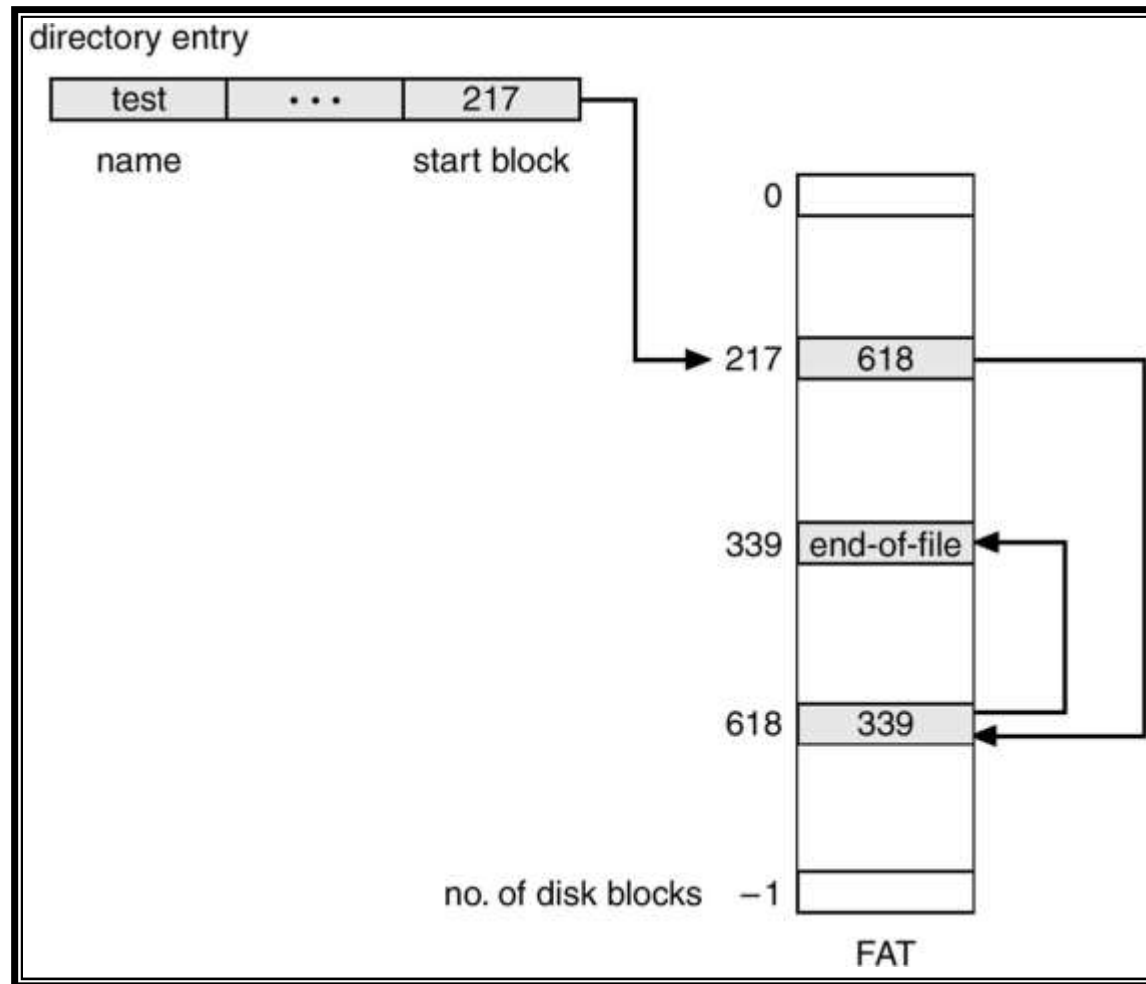


# *Linked Allocation*





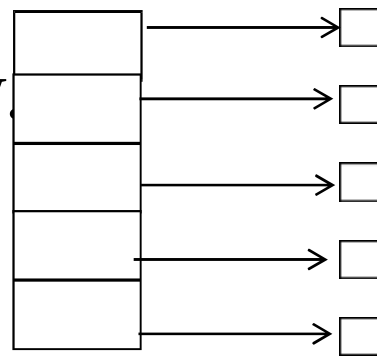
# *File-Allocation Table*





# *Indexed Allocation*

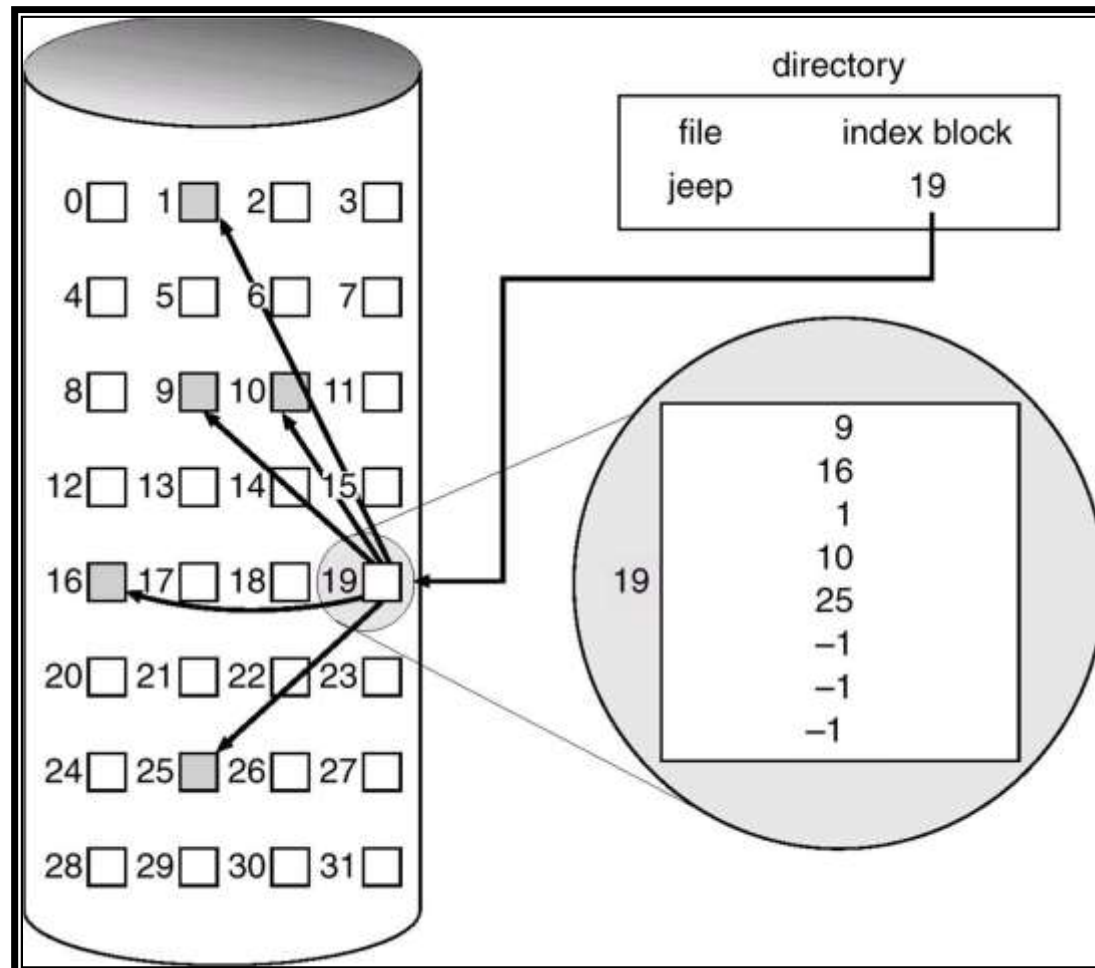
- Brings all pointers together into the *index block*.
- Logical view.



index table



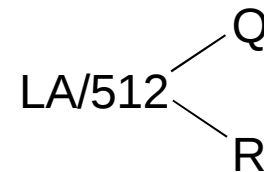
# *Example of Indexed Allocation*





## *Indexed Allocation (Cont.)*

- Need index table
- Random access
- Dynamic access without external fragmentation, but have overhead of index block.
- Mapping from logical to physical in a file of maximum size of 256K words and block size of 512 words. We need only 1 block for index table.



Q = displacement into index table

R = displacement into block



## *Indexed Allocation – Mapping (Cont.)*

- Mapping from logical to physical in a file of unbounded length (block size of 512 words).
- Linked scheme – Link blocks of index table (no limit on size).

$$LA / (512 \times 511) \begin{cases} Q_1 \\ R_1 \end{cases}$$

$Q_1$  = block of index table

$R_1$  is used as follows:

$$R_1 / 512 \begin{cases} Q_2 \\ R_2 \end{cases}$$

$Q_2$  = displacement into block of index table

$R_2$  displacement into block of file:



## ***Summarization***