

SNS COLLEGE OF TECHNOLOGY



Coimbatore-35
An Autonomous Institution

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DEPARTMENT OF ELECTRONICS & COMMUNICATION ENGINEERING

19ECT312 – EMBEDDED SYSTEM DESIGN

III YEAR/ VI SEMESTER

UNIT 4: Embedded Operating System and Modelling

TOPIC 4.8: POSIX Semaphores





- ➤ POSIX semaphores are synchronization primitives used in multi-threaded programming to control access to shared resources among concurrent threads
- ➤ Unlike mutexes, which allow only one thread to access a resource at a time, semaphores can permit multiple threads to access a resource simultaneously, up to a specified limit
- Semaphores maintain an internal counter that represents the number of available resources or permits, which threads acquire or release using the sem_wait() and sem_post() functions, respectively





This flexibility makes semaphores suitable for scenarios where multiple threads need controlled access to shared resources or where synchronization needs to be more granular than what mutexes offer

➤ However, improper usage of semaphores can lead to deadlocks or race conditions, so careful programming and understanding of concurrency principles are essential when working with POSIX semaphores.





POSIX Semaphors API

- ➤ POSIX (Portable Operating System Interface) semaphores API provides a standardized interface for controlling semaphores in Unix-like operating systems
- > Semaphores are synchronization primitives used for inter-process communication and coordination
- ➤ In POSIX, semaphores are typically used to coordinate access to shared resources among multiple processes or threads
- They can be thought of as counters with associated atomic operations for incrementing, decrementing, and testing their values

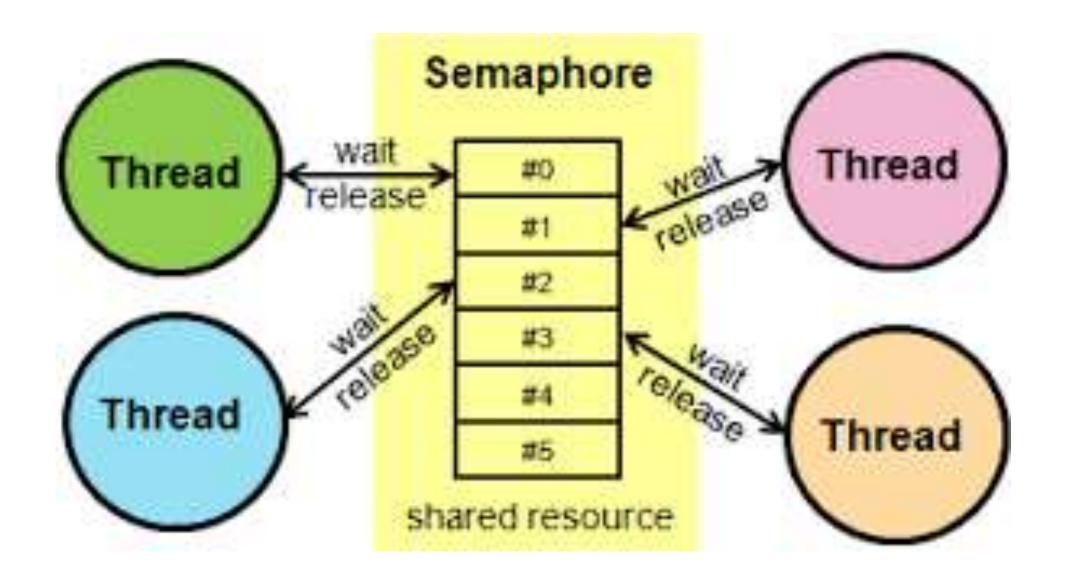




- 1. sem_init: Initializes a semaphore with a specified initial value.
- 2. sem_destroy: Destroys a semaphore, releasing any associated resources.
- **3. sem_wait:** Decrements the value of a semaphore. If the value is zero, the function blocks until the semaphore becomes non-zero.
- **4. sem_post**: Increments the value of a semaphore.
- **5. sem_getvalue**: Retrieves the current value of a semaphore without modifying it.











Advanced Semaphore Techniques

Advanced semaphore techniques involve more sophisticated usage patterns and scenarios beyond basic synchronization

Advanced techniques

1. Multiple Semaphores for Resource Allocation

- Instead of using a single semaphore to control access to a shared resource, you can use multiple semaphores to manage different aspects of resource allocation
- For example, one semaphore can control read access, another semaphore can control write access, and additional semaphores can manage other types of access or resource states.





2. Counting Semaphores

- ➤ While binary semaphores have only two states (0 and 1), counting semaphores can have an initial count greater than 1
- They are useful for scenarios where multiple instances of a resource can be allocated simultaneously
- Threads or processes decrement the semaphore count when they acquire the resource and increment it when they release it.





Advantage:

- ❖Portability: Standardized interface across Unix-like operating systems ensures compatibility and easy migration of code.
- ❖ Inter-Process Communication (IPC): Facilitates synchronization and communication between multiple processes.
- **Scalability**: Adaptable for simple to complex synchronization needs in applications with multiple processes or threads.
- **❖Flexibility**: Offers binary and counting semaphore types for diverse synchronization requirements.
- **Efficiency**: Implemented with efficient algorithms and system calls, minimizing overhead in memory and processing time.
- **Ease of Use**: Simple API with intuitive functions for semaphore management simplifies development and maintenance.





Limitations:

- **Limited Functionality**: Lack advanced features like deadlock detection and priority inheritance found in other synchronization primitives.
- **Complex Error Handling**: Error handling can be intricate, requiring careful attention to return values and error codes.
- **❖Kernel Dependency**: Performance and behavior may vary based on the underlying operating system and kernel version.
- ❖ Resource Overhead: Each semaphore consumes system resources, potentially becoming problematic in applications requiring many semaphores.
- ❖ Portability Challenges: While aiming for portability, differences in behavior and implementation across platforms may arise.
- *Risk of Deadlocks and Races: Improper use can lead to deadlocks or race conditions, demanding careful programming to avoid.





Thank you